

## Algorithmic Aspects of Communication – Assignments

The University of Tokyo – Winter 2013

Instructor: François Le Gall

Solve the three problems below. Submit your report to the instructor **before February 17th (5:00 PM)** by putting a copy in the instructor's mailbox located on the 1st floor of the Faculty of Science Building Number 7. Do not forget to write your name and your student number on the report.

### Problem 1 (Communication Complexity)

- (1) Let  $q$  be a prime number and let  $GF(q)$  denote the finite field of size  $q$ . Suppose that Alice has input  $x \in GF(q)$  encoded with  $\lceil \log_2 q \rceil$  bits and Bob has input  $y \in GF(q)$  encoded with  $\lceil \log_2 q \rceil$  bits. Let  $f_q: GF(q) \times GF(q) \rightarrow \{0, 1\}$  be the function defined as

$$f_q(x, y) = \begin{cases} 0 & \text{if } x + y + xy = 0 \\ 1 & \text{if } x + y + xy \neq 0 \end{cases}$$

for all  $x, y \in GF(q)$ . Show that the randomized communication complexity of  $f_q$  is  $O(\log \log q)$  bits. Can you prove a lower bound on the deterministic communication complexity of  $f_q$ ?

- (2) Let  $r$  be a positive integer and let  $V$  denote the vector space of dimension  $r$  over the finite field  $GF(2)$ . In other words,  $V = \{0, 1\}^r$ , where the addition of two strings is done component-wise modulo 2. Suppose that Alice has as input a subspace  $U \subseteq V$ , and Bob has a vector  $y \in V$ . The goal is for Bob to decide if  $y \in U$  or not. We suppose that only Alice can send messages to Bob, and that Alice wants to send as few communication as possible (this model is called *one-way communication complexity*). Construct a randomized communication protocol that solves (with high probability) this task using  $O(r)$  bits of communication. (Hint: consider the orthogonal subspace  $U^\perp$  of  $U$ .)

### Problem 2 (Private Information Retrieval)

In lecture 7, we discussed information-theoretic private information retrieval. In the case of a single server, it was mentioned that any protocol must exchange  $\Omega(n)$  bits of communication, where  $n$  is the size of the database owned by the server. Give a proof of this assertion.

### Problem 3 (Network Coding)

This problem deals with the  $k$ -pair problem in network coding. The graph of

Figure 1 has three sources  $s_1, s_2, s_3$  and three targets  $t_1, t_2, t_3$ . The goal is to send simultaneously one unit of information  $x$  (i.e., one element  $x \in GF(q)$  for some prime  $q$ ) from  $s_1$  to  $t_1$ , one unit of information  $y$  from  $s_2$  to  $t_2$  and one unit of information  $z$  from  $s_3$  to  $t_3$ . Each edge in the graph is again supposed to have unit capacity. We say that a protocol is linear if, for each edge  $(u, v)$ , the message sent through  $(u, v)$  is a linear combination (with coefficients in  $GF(q)$ ) of the messages received at node  $u$ .

- (1) Design a linear protocol solving this task, for  $q = 2$ .
- (2) Are there other prime numbers  $q$  for which a linear protocol exists?

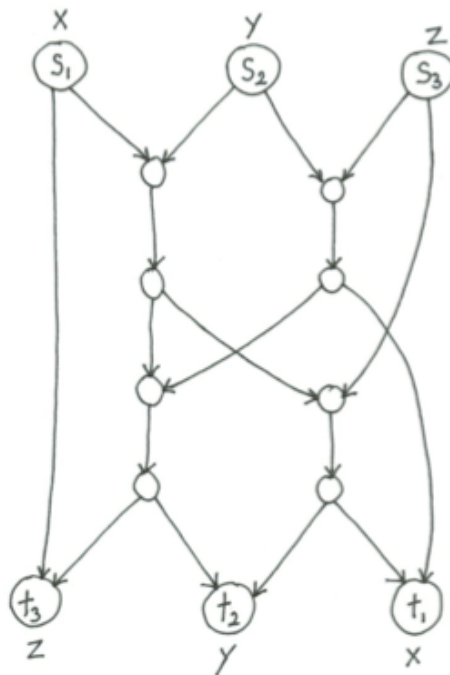


Figure 1: a three-source three-target graph